

# **A Modeling & Simulation Implementation Framework for Large-Scale Simulation**

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## **Abstract**

Classical High Level Architecture (HLA) systems are facing development problems for lack of supporting fine-grained component integration and interoperation in large-scale complex simulation applications. To provide efficient methods of this issue, an extensible, reusable and composable simulation framework is proposed. To promote the reusability from coarse-grained federate to fine-grained components, this paper proposes a modelling & simulation framework which consists of component-based architecture, modelling methods, and simulation services to support and simplify the process of complex simulation application construction. Moreover, a standard process and simulation tools are developed to ensure the rapid and effective development of simulation application.

**Keywords:** simulation environment, HLA federate, component, system architecture

## **1. Introduction**

The modeling and simulation technology is widely used in many fields such as industry, scientific research and military analysis. As the simulation applications become more and more complicated, constructing these systems become more difficult. There are measures to simplify the simulation development. Standard process models, such as Distributed Simulation Engineering and Execution Process (DSEEP) [1], define common, unified process models, which provide a common framework for different communities to describe their engineering practices, so the engineers can communicate with each other well.

And standard simulation interoperation protocols, such as HLA [2], DIS [3] and TENA [4] provide common distributed interoperation methods between different member applications, so subsystems can be composed to construct a more complex simulation system. This paper presents a simulation environment, which is intended to utilize the best practices in the field of simulation to support simulation applications for large-scale complex system.

The rest of this paper is organized as follows. Section 3 provides a brief introduction to the whole simulation development process and the architecture of the simulation environment. Simulation model components are discussed in Section 4, including the definition of components and how to compose these components. In Section 5, simulation services are described, which form the runtime for simulation models. At the last runtime, object database is presented in Section 6, which provides distributed communication capabilities.

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## 2. Related Work

Large-scale simulation requires advanced implementation simulation software approaches. Many pioneering works have been done to construct the theory base, including [1-6]. However, most of the works are focused on simulation process and components theoretic methods while they are lack of implementation techniques.

[11-12] proposed cell-DEVS and developed CD++ tools, which is open source and supporting discrete event system modeling and simulation. [13] designed HyperWarpSpeed time management algorithms, developed parallel simulation engine and used the engine in military simulation systems. [14] developed simulation engine GTW and the product named FDK, using in multi-field including biology, physiology, traffic and communication etc. All the works are useful for reference. However, they did not address the scheduling service management of local multi-components.

Moreover, HLA is also recognized with its low running efficiency because of losing fine-grained component integration and interoperation in local HLA federates. Here we have an underline prerequisite is that a lower HLA federation is made slow for the lowest federate waiting by other federates. Therefore, to provide more efficient methods of this issue, an extensible, reusable and compensable simulation framework is proposed in this paper and we start discussing the development of a promoted and standardized process and architecture.

## 3. Process and Architecture

A standard process can provides a common language for simulation developers to communicate and gives a better understanding to the life-cycle of the simulation application. To suit this component-based simulation environment, a process is proposed which tailored and specialized from DSEEP. The process is summarized as follows.

Step 1: Define simulation application objective. The user, the sponsor, and the development/integration team define and agree on a set of objectives and document what must be accomplished to achieve those objectives.

Step 2: Perform conceptual analysis. Scenario development and conceptual modeling take place in this step. The output including scenarios and conceptual models guides the rest of the development process.

Step 3: Design simulation application. Existing simulation model components which are reusable is selected, and new components are designed. Component model is described in the next section.

Step 4: Develop simulation application. The new simulation model components are implemented, and data exchange model is developed for Runtime Object Database.

Step 5: Integrate and test simulation application. Simulation model components are composed to form simulation entities, and simulation services suitable for this simulation are selected and integrated.

Step 6: Execute simulation and analyze results. The simulation is executed and output data from the execution is collected and analyzed.

The scenarios, conceptual models and model components created in the process can be stored in repositories for reuse in other simulation applications.

In the simulation environment, common simulation services are provided to form the simulation runtime, so the simulation application developers can focus on simulation modeling. The simulation environment consists of four layers, which are network communication layer, runtime object database layer, simulation services layer and models & tools layer.

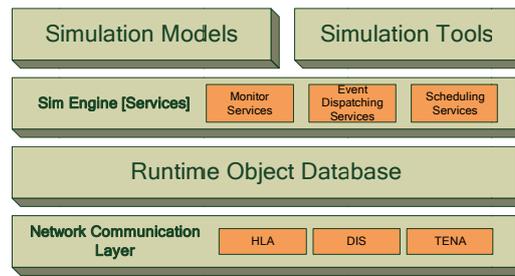


Fig. 1 Simulation environment architecture

#### 4. Simulation model components

Component-based composability is one of the key features of the simulation environment. Small granular models have lots of advantages compared to large granular models. For example, they are easy to understand and develop. So one of the best ways to simulate complex system is breaking the system into smaller sub models (also referred to as simulation model components), and then composing them together to form the whole system. Composability refers to the ability of a simulation to be flexible configured to adapt to a range of missions, scenarios and simulation models [7]. With composability, one can rapidly assemble and configure a unique simulation system from existing simulation components. These created simulation applications can vary in not only scenarios but also behaviors.

Conceptual model is a non-software description of the computer simulation model, describing the objective, inputs, outputs, contents, assumptions and simplifications of the model [8][9]. Although conceptual model is not a part of the component architecture, it is vital to model component development because it restrains what the components do and how they communicate. So the conceptual model formalism is briefly introduced here; and how it is converted to components is described in the following paragraphs.

In DEVS [6], a discrete event system is presented as

$$M = \langle X, S, Y, \delta_{int}, \delta_{ext}, \lambda, ta \rangle$$

where  $X$  is the set of input values;  $S$  is a set of states;  $Y$  is the set of output values;  $\delta_{int}: S \rightarrow S$  is the internal transition function;  $\delta_{ext}: Q \times X \rightarrow S$  is the external transition function, where  $Q = \{(s, e) | s \in S, 0 \leq e \leq ta(s)\}$  is the total state set,  $e$  is the time elapsed since last transition;  $\lambda: S \rightarrow Y$  is the output function;  $ta: S \rightarrow R_{0, \infty}^+$  is the set positive reals with 0 and  $\infty$ . In the simulation environment,  $X$  and  $Y$  are mapped to interfaces and events of components,  $S$  are mapped to properties of components, and  $\delta_{int}$ ,  $\delta_{ext}$ ,  $\lambda$ ,  $ta$  are translated to implementations of components.

##### 4.1. Interfaces

An interface declares a set of operations that a client may request. A component which satisfies the interface provides its service through the operations of this interface. Components communicate only through interfaces they implement, so it is easy to reuse and extend components as the whole system is described by interfaces not specific components.

##### 4.2. Events

The notion of events is also described by interfaces. There are two specific interfaces which are EventSource and EventSink. Components which support events must implement one or both of them. An event source embodies the potential for the component to generate events of a specified type. An event sink embodies the potential for the component to receive events of a specified type.

```
interface EventSource {
    Events getEvents();
}
interface EventSink {
    void pushEvent(in Event e);
}
```

The event dispatch strategy is not specified by the components but decided when composing the components into a system. Two dispatch strategies are provided by now, and more can be added. To solve the problem of confliction, each type of events can only has one dispatch strategy. The publish/subscribe strategy is that the event sources, sinks publish, subscribe types of events, and event dispatching service transmits the specified types of events from publishers to subscribers. The connect strategy is defined by connections between components with each connection linking two components. Events are transmitted through the connections.

#### 4.3. Components

There are two types of DEVS components: atomic components and composed components [6]. A composed component consists of atomic components and/or other composed components, while atomic components cannot be divided into smaller parts.

The definition of a component includes interfaces, references, input events and output events. Interfaces indicate the operations provided by this component. References are connection points to other components, so this component can use services supplied by the referred object. Input and output events are the types of events used to communicate with other components.

#### 4.4. Component Composition

Component composition is implemented by references which link to subcomponents. References to other components are usually though not specific components but interfaces. The component can refer to any component which satisfies the specified interface. Component connections are specified by an assembling script based on XML. The script describes which instances of components are in the system, their initialization parameters and how they connect to each other.

In this simulation environment, simulation entity is a kind of component, which satisfies special constrains including the initial parameter loading interface and simulation engine event interfaces. Simulation entities are relatively independent components, because they do not have references to other components and they communicate only through events. Simulation entities are the smallest process units scheduled by the simulation runtime.

## 5. Simulation services

### 5.1. Simulation Engine

Simulation services forms Simulation engine, which includes simulation monitor service, time management service, event dispatching service and simulation scheduling service. The engine drives simulation models to advance the simulation and generate output data. The kernel classes of the simulation engine are shown in Fig. 2.

Entity and Event are the fundamental elements of discrete event simulation. Events are passed form Event Sources to Event Sinks by Event Dispatcher. Scheduler manages Logical Processes which manipulate the runtime states of the Entities. The running process is illustrated in Fig. 3 .

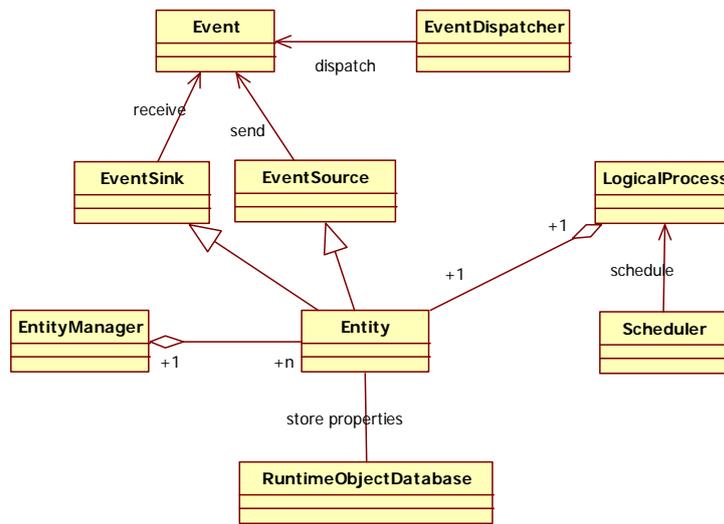


Fig. 2 Kernel classes

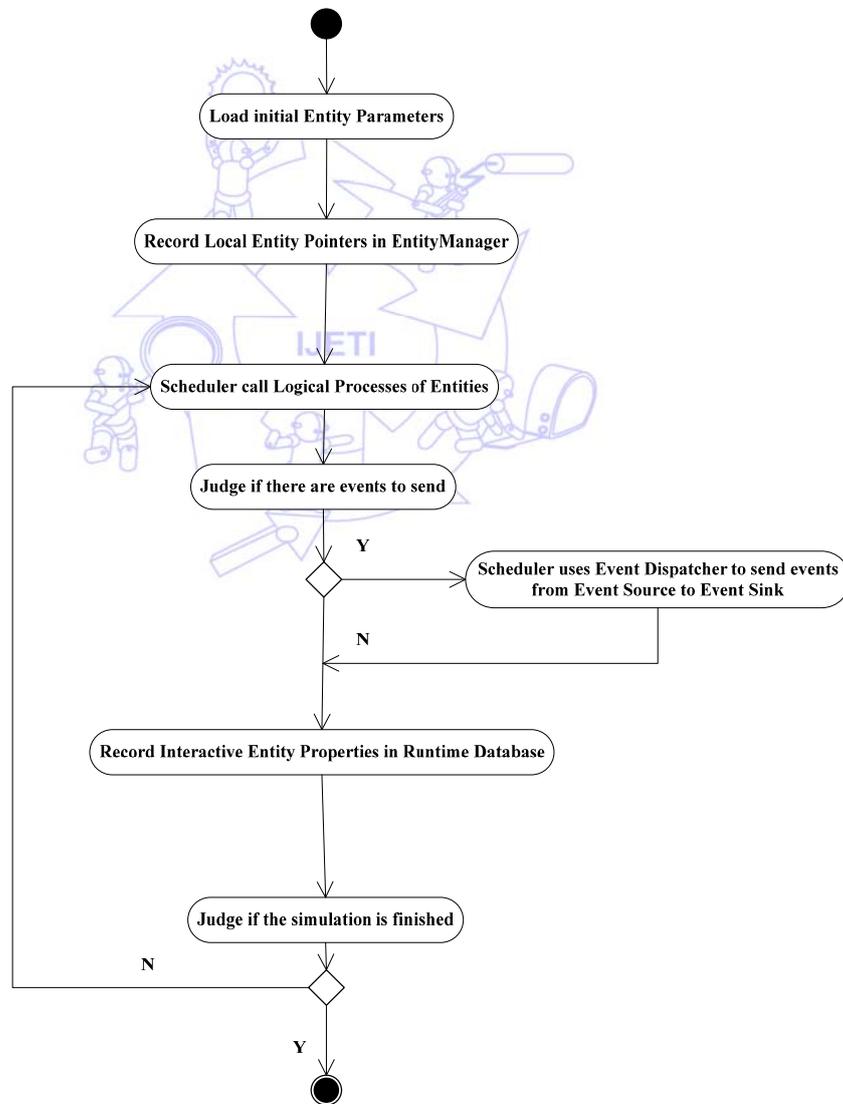


Fig. 3 Running Process of Simulation Engine

### 5.2. Simulation Monitor Service

Simulation monitor service controls the working state of the simulation engine. These states include the initializing, suspending, running, saving, restoring and terminating states. After the simulation engine has been initialized, it goes automatically into the suspending state. In the simulation process, the engine can be in suspending or running state, controlled by the commands from the client. When the simulation terminating condition is satisfied or the monitor receives terminating commands, the simulation engine goes into terminating state, and the simulation process finishes.

### 5.3. Event Dispatching Service

The notion of events has been introduced above, and two event dispatch strategies between components have been proposed. At the inter-entity level, a third strategy is provided to enrich the semantics of events. As each simulation entity has a unique identifier, so simulation entities can send events to the certain entities which they want to communicate with by using this unique identifier. The types of events using this dispatch strategy must have an identifier to indicate the event target and a time tag, and the events are delivered to the target specified by target identifier and sorted by the time tag. More strategies can be added at the inter-entity level, and events must register their dispatch strategies before simulation starts.

### 5.4. Scheduling Service

Scheduling service is responsible for scheduling simulation models to processing units. There are two ways which are serialization and parallelization to do this. Serialization way is relatively simpler and runs faster on single processing unit, while parallelization way is more effective on multiple processing units, because it can utilize the parallel computation capability of multi-processors.

The simulation entities are processed one by one in the serialization method, which is described as follows.

```
for each simulation entity  $Entity_i$  {
  for each simulation event  $Event_{ij}$  of  $Entity_i$  {
    process  $Event_{ij}$ 
  }
}
```

The parallelization method consists of a **monitor thread** and multiple **worker threads**. The monitor thread is responsible for computation task assignment to worker threads, while the worker threads complete the computation tasks. The scheduling is committed at the simulation entity level, which means a simulation entity with all its child components is assigned to one worker thread. As simulation entities do not refer to other simulation entities, the worker threads do not have shared variables and the simulation can be paralleled without synchronization between worker threads. The algorithm is presented below.

```
Worker thread:
while the simulation is running {
  if there are computation task  $Task_k$  {
    run  $Task_k$ 
  } else
  notify the monitor thread and wait for next task
}
Monitor thread:
while there are simulation entities to be processed {
  if worker thread  $Thread_i$  is idle {
    assign one of the simulation entities to  $Thread_i$ 
  }
}
```

### 5.5. Event-driven Conservative Time Synchronization Mechanism

An important concept in the field of distributed simulations is time synchronization mechanism, which maintains a consistent causality in the process. This paper uses the event-driven conservative time management mechanism.

*Monitor thread* receives the Lower Bound on Time Stamp (LBTS) in the message format from the *Worker threads*, calculates the minimum of all LBTSs, and then returns the minimum as Greatest Available Logical Time (GALT) to all the *Worker threads*.

According to logical time management strategy in HLA, the relationship of a federate to other federates is divided into two types: logical time regulating and logical time constrained. So there are four logical time management states: not only logical time-regulating but also constrained, neither logical time-regulating nor constrained, only logical time-regulating and only logical time constrained. Accordingly, *Worker threads* are divided into four types in this paper:

Only logical time regulating *Worker threads*: Its time synchronization affects other *Worker threads*' time synchronization without being affected by the others. So its LBTS is constantly 1. Only logical time constrained *Worker threads*: Its time synchronization is affected by other *Worker threads*' time synchronization without affecting others' time synchronization. So its LBTS can be calculated by the formula below:

$$LBTS_i = \text{Min}(Ltime_j + Lookhead_j)$$

$Ltime_j$  is the current logical time of *Worker thread j*,  $Lookhead_j$  is the looking ahead time of *Worker thread j*,  $i, j = 1, 2, \dots, n$ .

Not only logical time regulating but also constrained *Worker thread*: Its time synchronization not only affects other *Worker thread*'s time synchronization but also is affected by other *Worker thread*'s time synchronization. When it is considered as the logical time regulating *Worker thread*, its LBTS is 1. When it is considered as the logical time constrained *Worker thread*, its LBTS can be calculated according to only logical time constrained *Worker thread*. Neither logical time regulating nor constrained *Worker thread*; It is not referred in this platform, so we do not discuss it in this paper.

## 6. Runtime object database

Runtime object database is a virtual database, which provides a unified method to represent simulation object data and the transparent distributed communication capabilities. The database supports publish/subscribe facility. Simulation entities publish their states to the database and subscribe the messages they interested in. Runtime object database supplies the interoperation environment between simulation entities, and provides data access ports for simulation tools, such as visualization tools and data collection tools. It hides network communication protocols from the simulation models and simulation services above this database. The relationship between runtime object database and other modules is shown below.

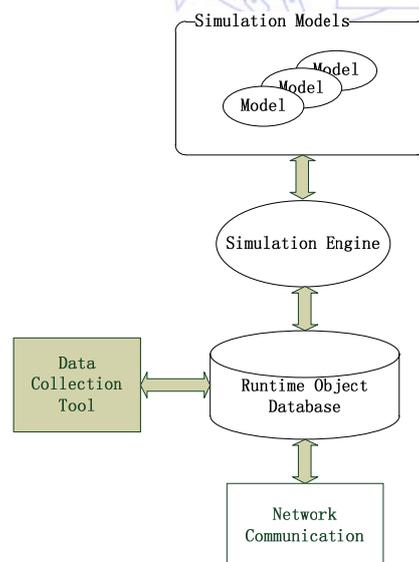


Fig. 5 Runtime object database

Runtime object data model is defined by simulation objects and simulation interactions. Simulation objects which are described by properties map to simulation entities, while interactions represented by parameters map to simulation events. All of the data saved into this database must satisfy the runtime object data model. There are a layer between the simulation models

and the database which translate the model data into runtime object data model, and a layer between the database and the network communication service which translate the runtime object data model to data structures of the specific network communication protocol.

## 7. Conclusion

This paper presents the architecture of a modeling and simulation environment, in which simulation components are composed to construct the simulation application. This simulation environment provides common simulation services and offers every chance to reuse existed resources to simplify development of complex simulation systems. By using of the runtime object database and software adapters including HLA-DEVS and DEVS-DIS agents, this environment can be compatible to many distributed interoperation protocols, such as HLA, DIS et cetera. By composing simulation services, this simulation environment can be tailored and specified to meet the needs of certain simulation applications.

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